



# TC

INSTITUT NATIONAL DES SCIENCES APPLIQUÉES DE LYON

## IST Semester

## Networks - Part 1 -

# CSMA/CA and IEEE 802.11



# CSMA/CA and IEEE 802.11

- **Speaker**

- Razvan Stanica
- <http://perso.citi.insa-lyon.fr/rstanica/IST-NET-WiFi.pdf>

- **Lab**

- CITI [*Centre of Innovation in Telecommunications and Integration of services*]
- INRIA UrbaNet team

- **Research**

- Mobile networking
- Medium access control in wireless networks
- Intelligent transportation systems





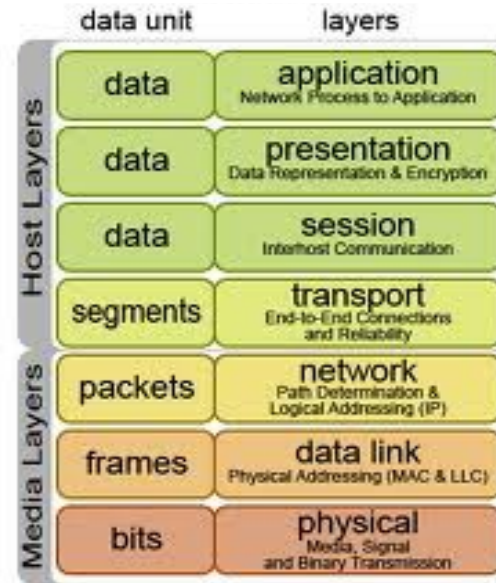
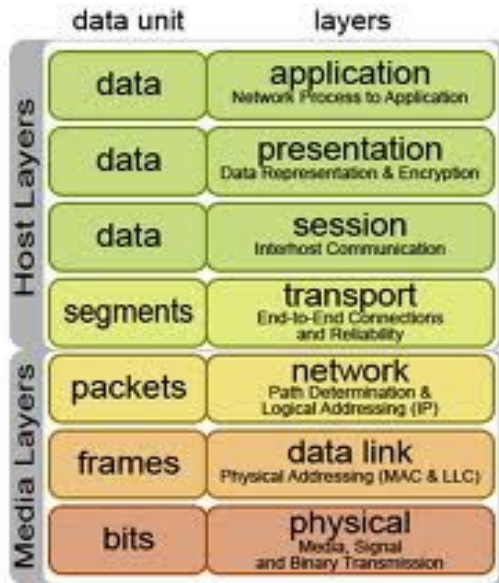
# CSMA/CA and IEEE 802.11

- The plan for today
  - What makes wireless so special?
  - Carrier Sense Multiple Access with Collision Avoidance – Why? How?
  - What else is in IEEE 802.11?





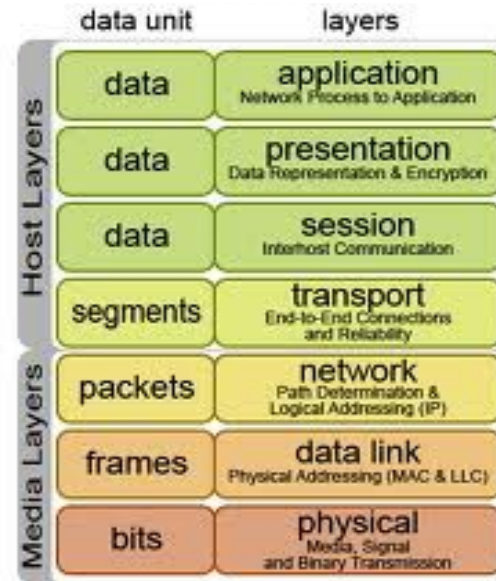
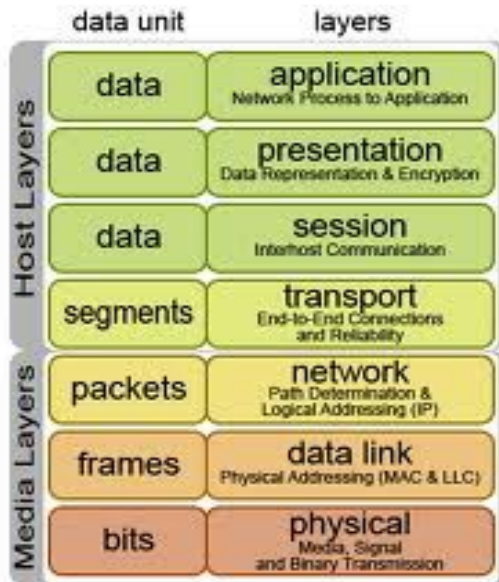
# Basic concepts: Networking







# Basic concepts: Networking



Wireless simply means that there are no wires at the PHY layer



# Basic concepts: PHY Layer

- Performed Services

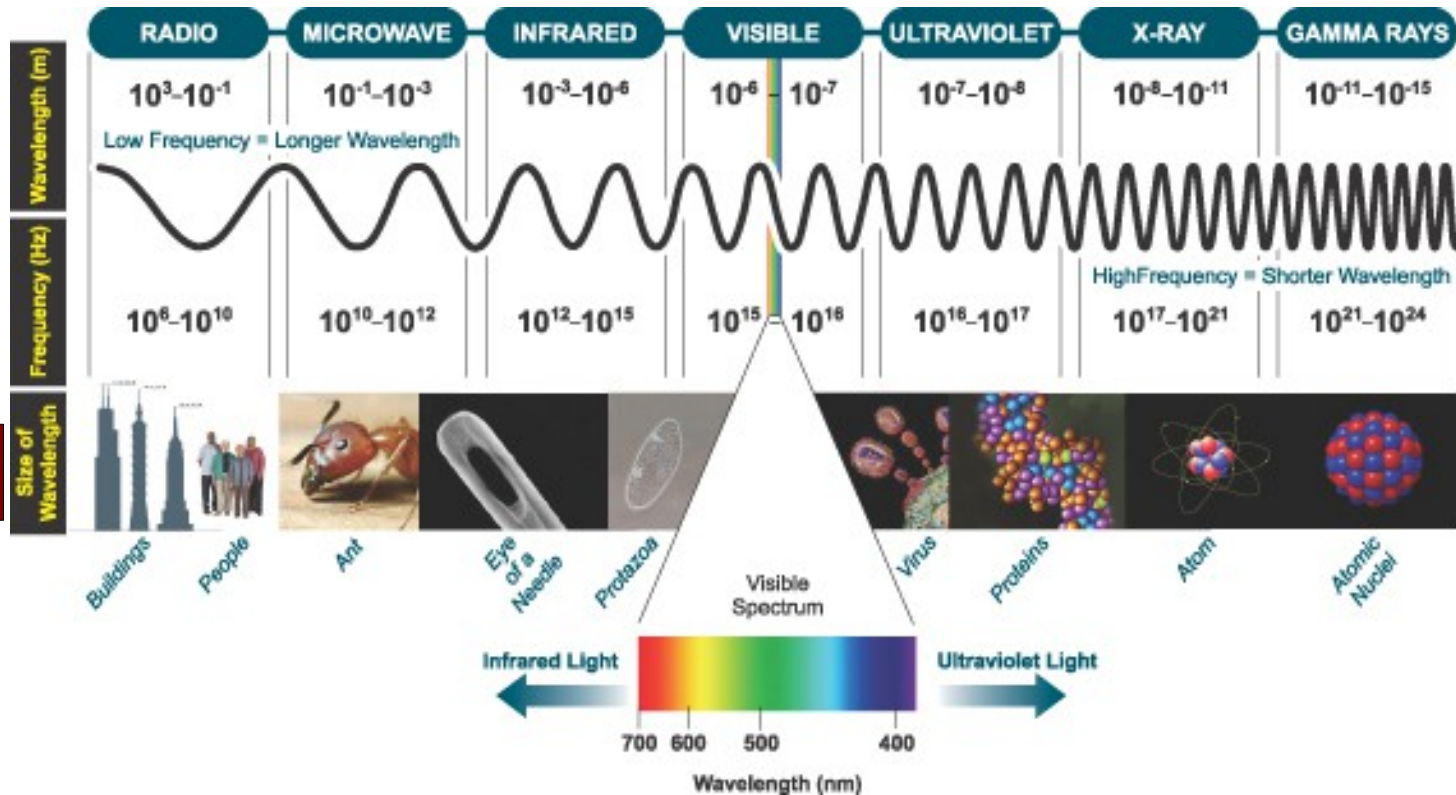


- Standardized mechanical and electrical interfaces
- Signal processing (equalization, training sequences)
- Coding (channel, line)
- Modulation (adding information to the EM wave)
- Carrier sense, collision detection (when possible)



# Basic concepts: Wireless

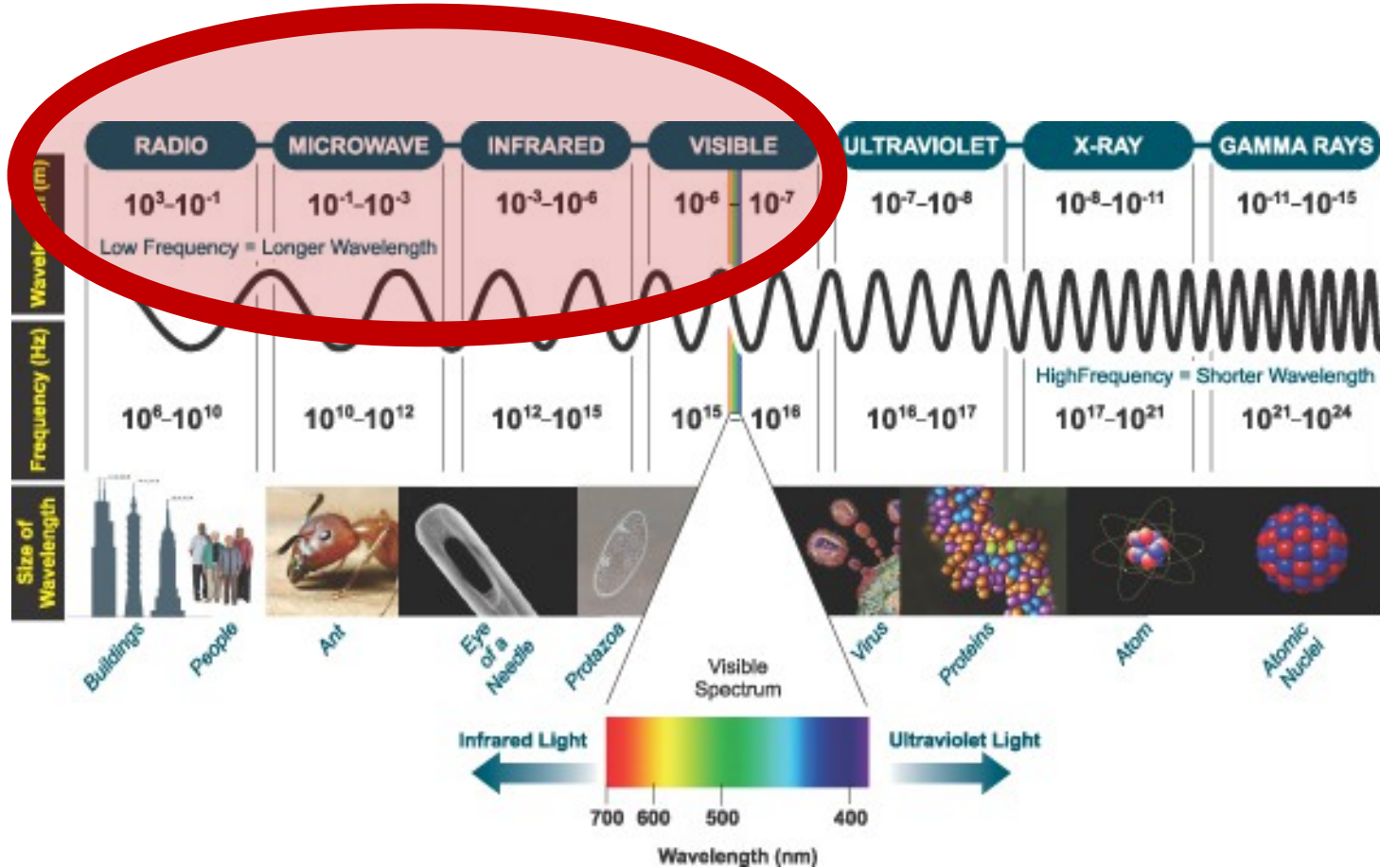
- EM frequency spectrum





# Basic concepts: Wireless

- Where wireless communication happens







# Basic concepts: Wireless Spectrum

Name	f	$\lambda$	Usage
Low Frequency	30 KHz	10 Km	Aeronautical & maritime navigation, metereology
Medium Frequency	300 KHz	1 Km	AM radio
High Frequency	3 MHz	100 m	Amateur radio (Morse code), marine, aviation, military
Very High Frequency	30 MHz	10 m	FM radio
Ultra High Frequency	300 MHz	1 m	Television, Cellular networks, WiFi
Super High Frequency	3 GHz	10 cm	WiFi, Satellite transmission, Bluetooth, Wireless USB
Extremely High Frequency	30 GHz	1 cm	Radar, Satellite sensing, Wireless HD, WiFi(?)
Infrared	300 GHz	1 mm	Lasers, LEDs, Free Space Optical Communication

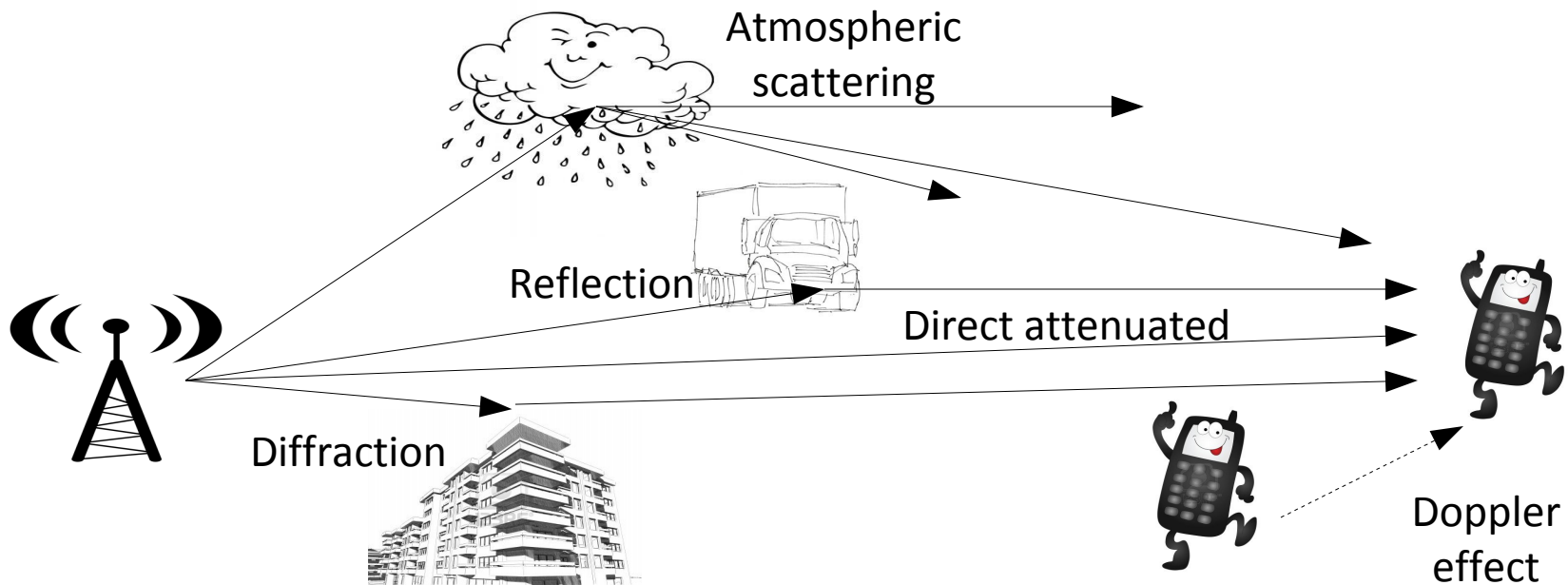




# Basic concepts: Wireless Propagation

- Everything you know from EM physics

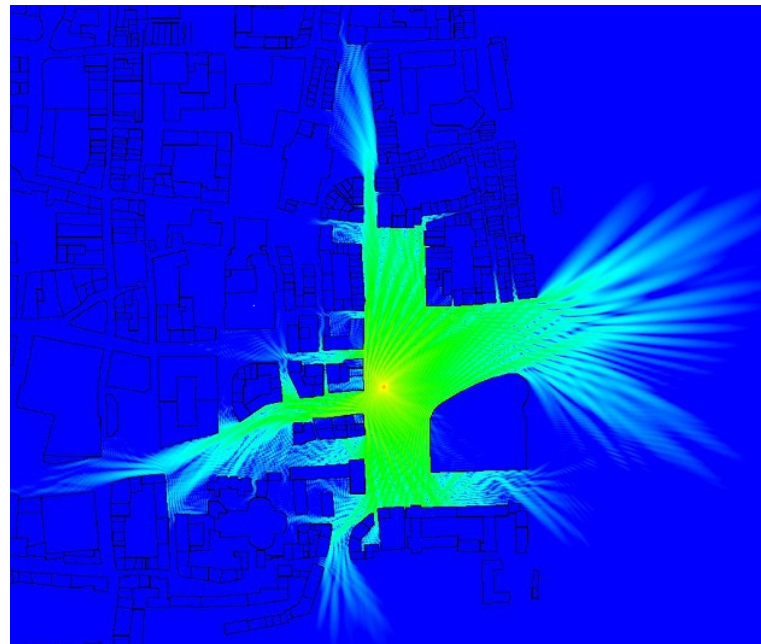
- Free space attenuation  $P_r \sim P_t \left( \frac{\lambda}{R} \right)^2$





# Basic concepts: Wireless Propagation

- As a result of the complex RF propagation, the **strength** of the received signal is highly irregular

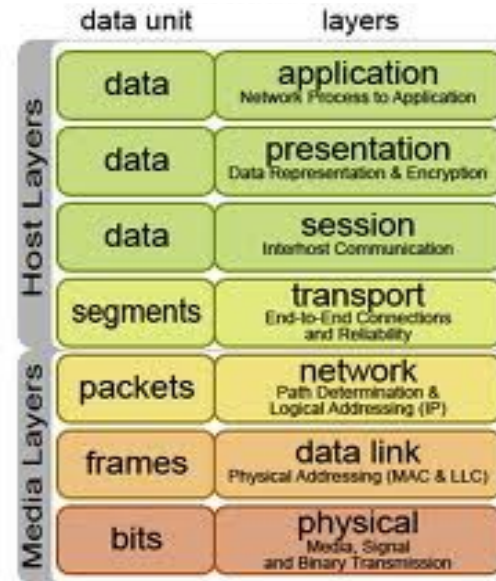
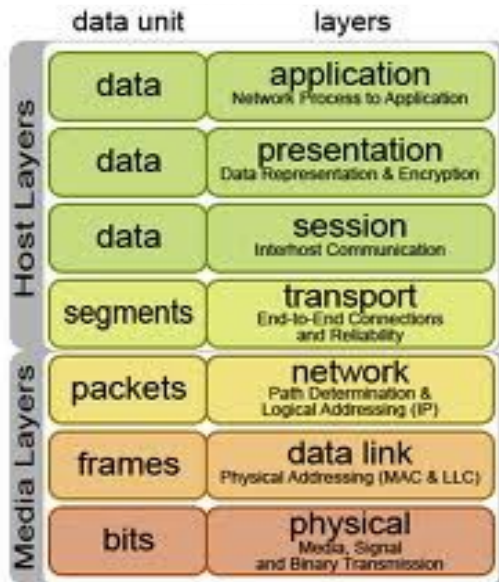


[ City of Koln, 100mW at 2.4GHz ]





# Basic concepts: MAC layer



The topic of this class is medium access control (a part of the data link layer)





# Basic concepts: MAC layer

- Why do we need a course for a sub-layer?
  - The MAC protocol decides who transmits at a given time and space
  - The number of different protocols is impressive: TDMA, CDMA, OFDMA, Token Bus, Token Ring, Aloha, etc.
  - This class is focused on Carrier Sense Multiple Access (but not 100%)





# CSMA: The principle

- Listen before you talk
  - Every station senses the channel for a certain time before transmitting
  - If the channel is idle, transmission goes on
  - If the channel is busy, wait and:
    - send as soon as it becomes idle (1-persistent)
    - choose a random back-off and try again (non-persistent)
    - send with probability  $p$  when idle ( $p$ -persistent)





# CSMA: The flavours

- Two major branches of protocols
  - CSMA with Collision Detection (CSMA/CD).  
You saw it (this morning?) in Ethernet.
  - CSMA with Collision Avoidance (CSMA/CA).  
Used in WiFi and ZigBee.





# CSMA/CD: A remainder

- How does it work?
  - It's CSMA, so we listen before we talk.
  - The CD part explains what happens when two stations start transmitting at the same time.
  - The collision detected procedure:
    - keep transmitting a jamming signal for one slot. This way everybody detects the collision.
    - check if you are allowed to retransmit (retransmission limit).
    - double the contention window and back-off.



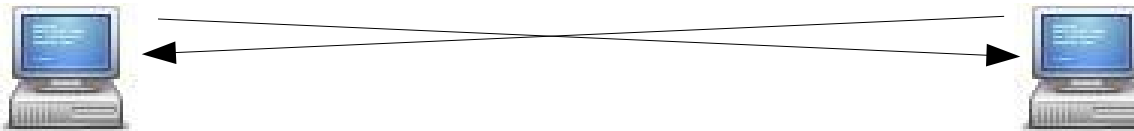
$b = \text{rand}(0, CW-1)$   
 $CW_{\min} = 1$





# CSMA/CD: A remainder

- How does it work?



signal power

- Ethernet uses a wired PHY
- Small signal attenuation
- Collision is detected from the first slot ( $5.12 \mu\text{s}$ )





# CSMA/CD: A remainder

- How does it work?

Station A	i					
Station B	i					
Station C	i					





# CSMA/CD: A remainder

- How does it work?

Station A	i	x				
Station B	i	x				
Station C	i					





# CSMA/CD: A remainder

- How does it work?

Station A	i	x				
Station B	i	x				
Station C	i					

- 4 possibilities for the back-off:
  - $b(A)=0, b(B)=0$
  - $b(A)=0, b(B)=1$
  - $b(A)=1, b(B)=0$
  - $b(A)=1, b(B)=1$







# CSMA/CD: A remainder

- How does it work?

Station A	i	x	x			
Station B	i	x				
Station C	i					

- On Ethernet, a small back-off can solve collisions quickly with high probability

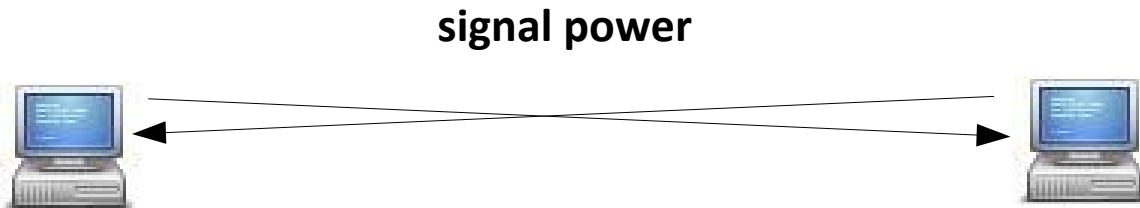




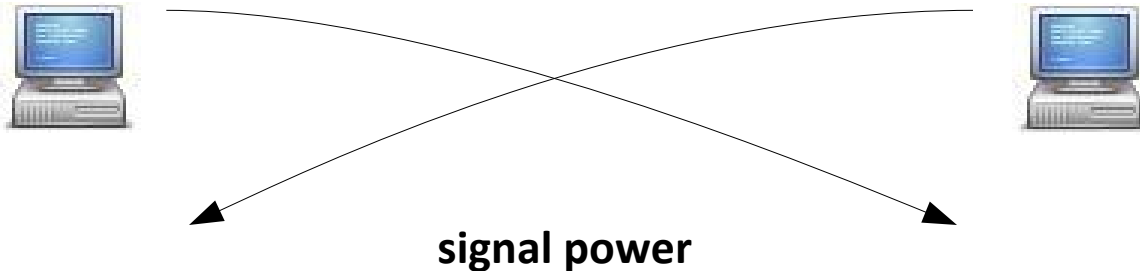
# CSMA/CD: A remainder

- Works great on Ethernet. Why not use it on WiFi?

- Wired



- Wireless – no CD while transmitting!



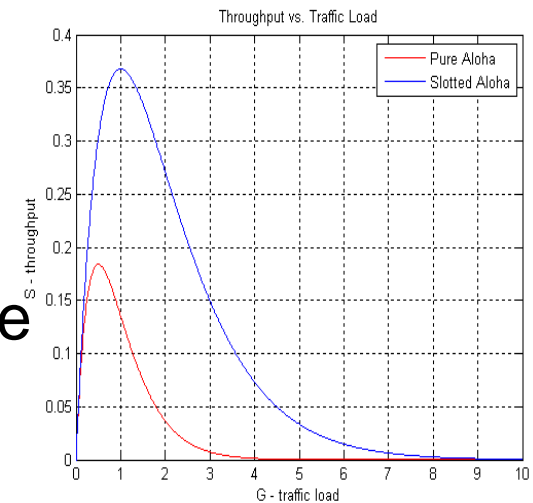
Signal attenuates much faster: The near-far effect





# CSMA/CA: The origins

- Wireless needs something different
  - Aloha
    - AlohaNet: first wireless packet network ('70s)
    - no carrier sense, just transmit
    - no ACK message means back-off
  - Slotted Aloha
    - stations are synchronized
    - transmissions can only start at the beginning of a slot





# CSMA/CA: The origins

- Wireless needs something different
  - Aloha is not efficient – too many collisions as the load increases.
  - The idea of using missing ACKs to schedule retransmissions is valid (and used in CSMA/CA).
  - However, missing ACKs **are not collision detection**. They are just a mechanism for error control – Automatic Repeat reQuest (ARQ)





# CSMA/CA: The origins

- Wireless needs something different
  - CSMA/CA appears in the '80s, originally as a MAC protocol in the Apple LocalTalk network.
  - It becomes the basis of the IEEE 802.11 standard, a.k.a. WiFi.
  - It is the most successful wireless communications technology up to date.







# CSMA/CA: The principles

- A station that wants to transmit

- Listen the channel (carrier sense)
- If channel is idle, transmit
- If channel is busy, choose a random back-off
- During back-off, time becomes slotted
- If a slot is idle, decrement the back-off timer
- If a slot is busy, freeze the back-off timer
- When the timer reaches 0, transmit
- If no ACK message is received, double the contention window and restart the procedure

$b = \text{rand}(0, CW-1)$

$CW_{\min} = 32$





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# CSMA/CA: The principles

- Why does it work?
  - The relatively high value for CW is essential.

Station A	i				
Station B	i				
Station C					





# CSMA/CA: The principles

i= idle, t= transmission,  
W= wait for ACK,  
x= collision detected

- Why does it work?

- The relatively high value for CW is essential.

Station A	i	t	t	t	t	t	w	w	x		
Station B	i	t	t	t	t	t	w	w	x		
Station C											

- On Ethernet, a collision was detected in maximum 5.12  $\mu$ s
- On WiFi, a missing ACK (not necessarily a collision), is detected after several ms





# CSMA/CA: The principles

i= idle, t= transmission,  
w= wait for ACK,  
x= collision detected

- Why does it work?

- The relatively high value for CW is essential.

Station A	i	t	t	t	t	t	w	w	x		
Station B	i	t	t	t	t	t	w	w	x		
Station C											

- Using  $CW_{min}$ , on CSMA/CD we have a 50% probability for another collision.
- Using  $CW_{min}$ , on CSMA/CA we have a 3,125% probability for another collision.



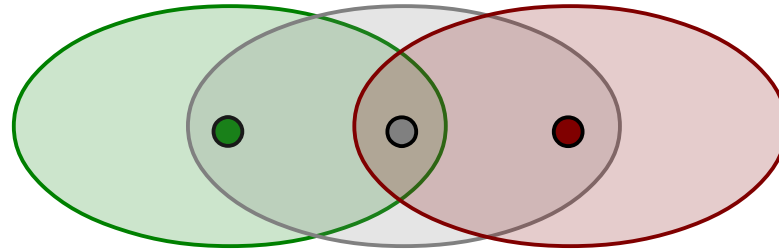




# CSMA/CA: The problems

- Great, this solves everything!
- Not exactly: hidden terminals

● and ● cannot hear each other



## Common scenarios

- sends DATA to ●
- wants to access the channel but hears the DATA from ●
- waits the end of DATA, then senses the channel idle and transmits
- ▶ collision between DATA and ACK

- sends DATA to ●
- wants to access the channel and senses it idle and transmits
- ▶ collision between DATA and DATA





# CSMA/CA: The problems

- The RTS/CTS handshake
  - The collision probability between hidden nodes increases with the size of the messages
  - The idea: use two short control messages to reserve the medium: Request-to-Send (RTS) and Clear-to-Send (CTS)
  - RTS and CTS contain information about the duration of the following transmission (DATA+ACK)

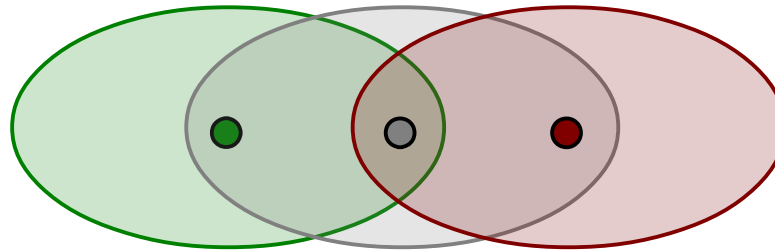




# CSMA/CA: The problems

- The RTS/CTS handshake

● and ● cannot hear each other



- sends RTS to ●
- sends CTS to ●
- hears the CTS and considers the channel busy for the announced duration
- sends DATA to ●
- sends ACK to ●

The RTS/CTS handshake is also known as virtual carrier sense.





# IEEE 802.11: The beginnings



- In 1985, the US Federal Communications Commission (FCC) created the Industrial, Scientific and Medical band (ISM) for non-licensed applications (2,4GHz)
- In 1990 the IEEE establishes the 802.11 committee
- The IEEE 802.11 standard was finalized in 1997 and became the de-facto standard for WLAN



# IEEE 802.11: Evolutions



- Higher data rate: 11 Mbps (b), 54 Mbps (g), 100+ Mbps (n)  
500+ Mbps (ac)
- Use of different frequencies (a – 5GHz, ad – 60GHz)
- Use of multiple antennas (n, ac)
- Integrating Quality of Service (e)
- Dedicated environments: mesh (s), vehicular (p)
- Security enhancements (i)

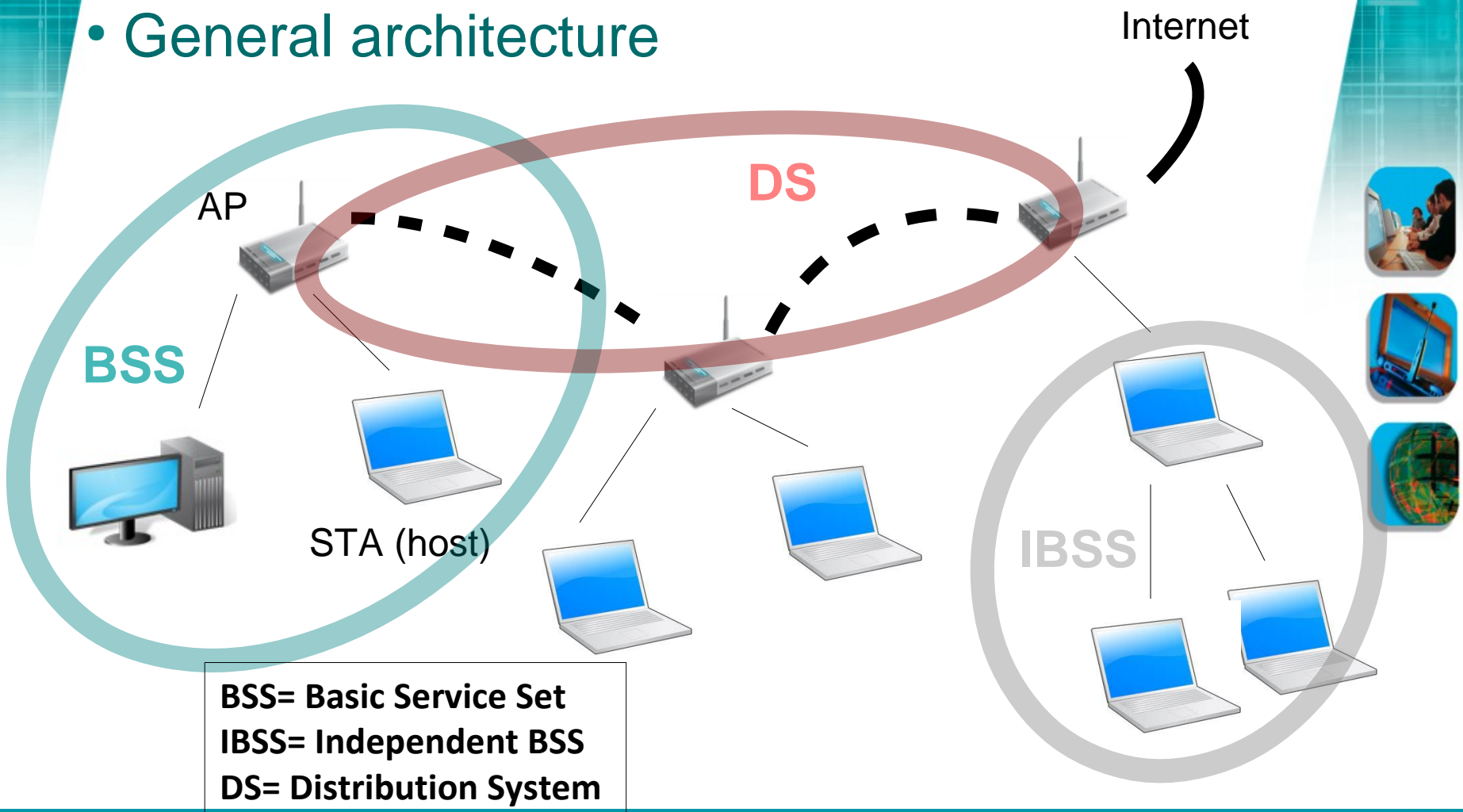






# IEEE 802.11: Introduction

- General architecture





# IEEE 802.11: Introduction

## • BSS association

- **Scanning** – STA looks for (chooses) an AP nearby
  - *Passive*: just wait for the periodic AP beacon
  - *Active*: probe APs for beacons
- **Authentication** – STA proves to have access
  - *Open*: skip this phase
  - *Secure*: challenge by the AP, the STA needs to have a shared key to respond correctly
- **Association** – STA enters the BSS
  - STA -> AP: association request
  - AP -> STA: association response
  - AP informs old AP via the DS in case of roaming



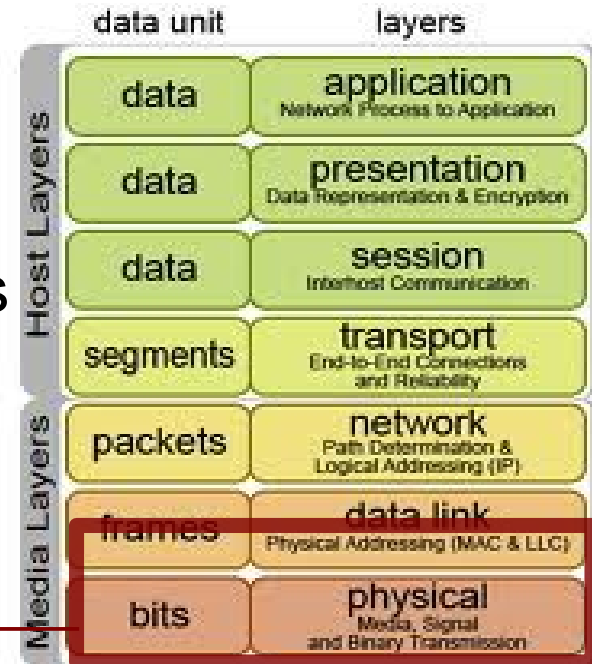


# IEEE 802.11: Introduction

- 802.11 is not only about the MAC

- IEEE 802.11 defines protocols for the **PHY** and **MAC**

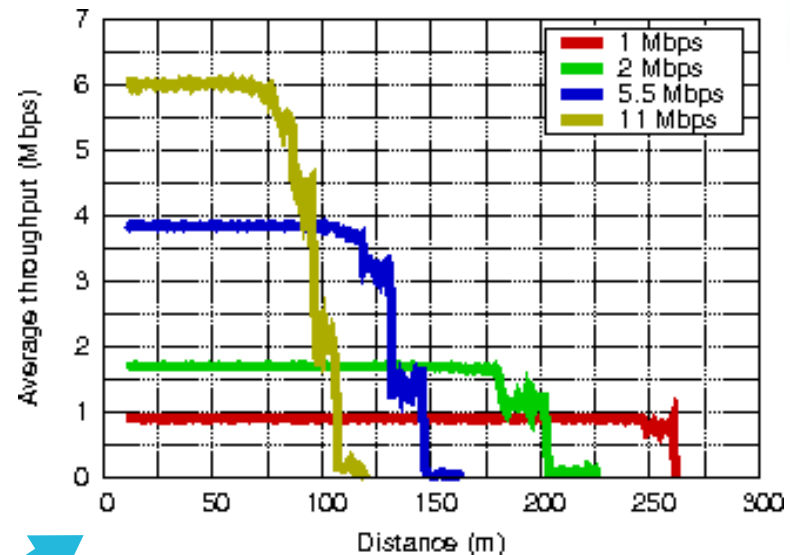
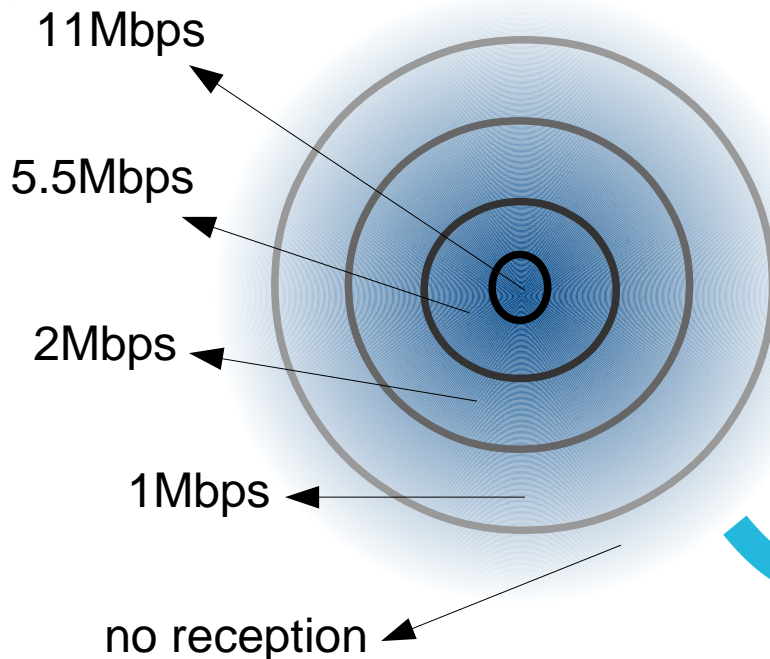
physical layer      medium access control sub-layer





# IEEE 802.11: PHY Basics

- Modulation – gives us the data rate
  - To decode a more complex modulation, we need a higher signal strength
  - Trade-off between data rate and transmission range



[ ns-2 simulation, 802.11b ]





# IEEE 802.11: DCF

- CSMA/CA implementation = Distributed Coordination Function (DCF)
  - four types of **InterFrame Space** (IFS)
    - Short InterFrame Space (**SIFS**)  
*used to separate transmissions belonging to a same dialogue (before a CTS or an ACK)*
    - Point coordination InterFrame Space (**PIFS**)  
*for data in the contention-free period (see later), to preempt any contention-based traffic*
    - Distributed InterFrame Space (**DIFS**)  
*standard IFS, used to separate transmissions of different dialogues*
    - Extended InterFrame Space (**EIFS**)  
*used by a station that received an erroneous frame*
- $SIFS < PIFS < DIFS < EIFS$

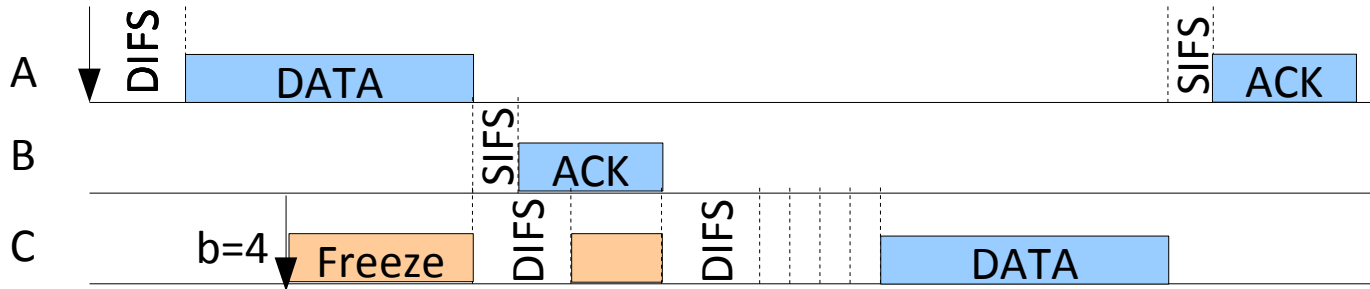






# IEEE 802.11: DCF

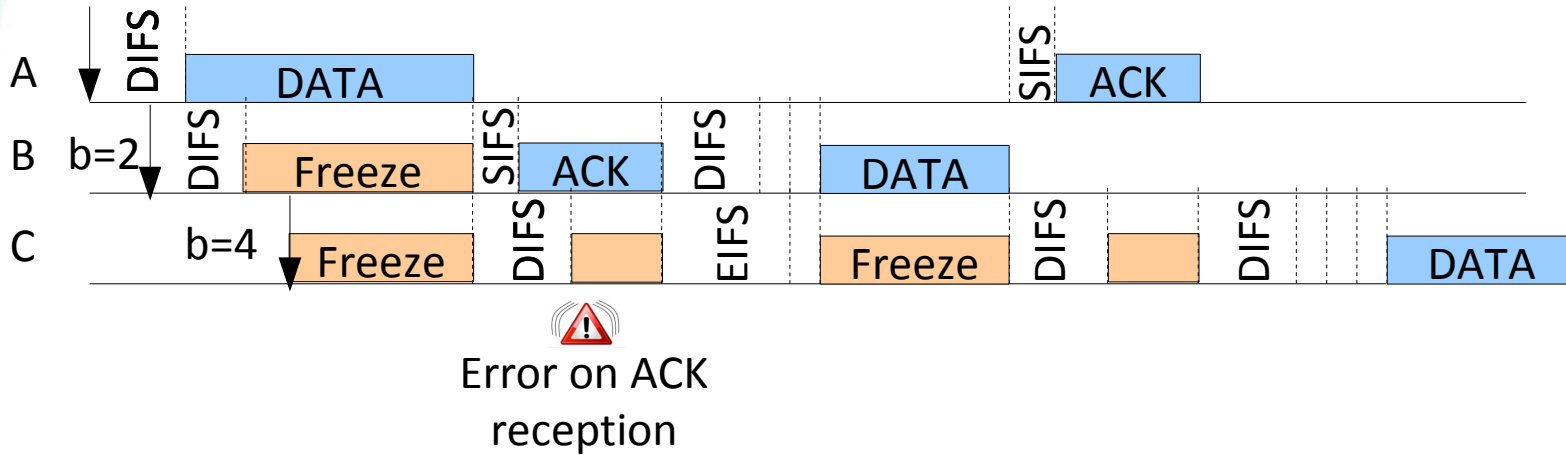
- Scenarios





# IEEE 802.11: DCF

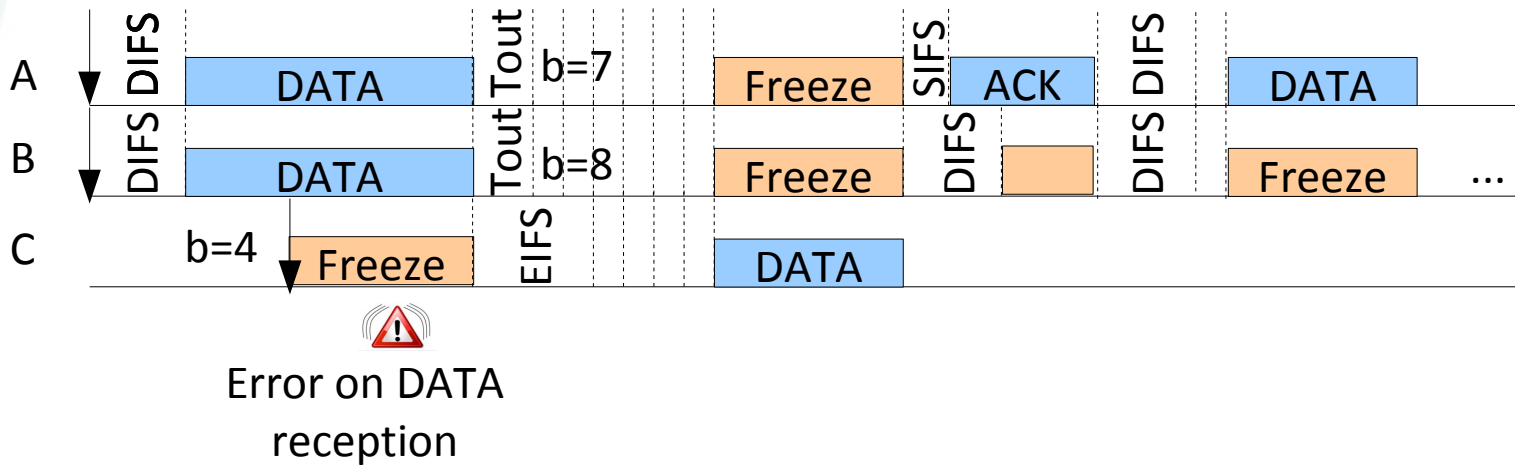
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# IEEE 802.11: DCF

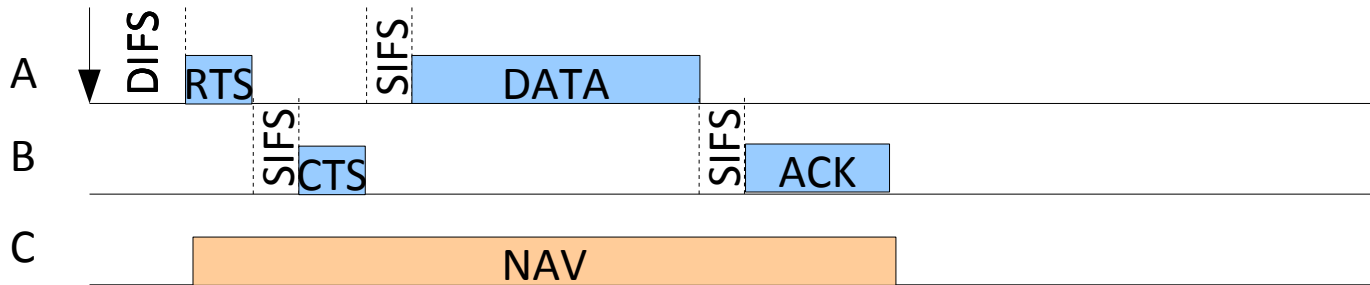
- Scenarios





# IEEE 802.11: DCF

- Scenarios



- NAV = Network Allocation Vector = Virtual Carrier Sense
- The RTS/CTS handshake is optional in IEEE 802.11





# IEEE 802.11: DCF

## Broadcast messages

- Broadcast= one transmitter, multiple receivers
- If all receivers transmit CTS or ACK after SIFS, collisions are unavoidable
- Broadcast messages are transmitted only once using the minimum CW, and their transmission is unreliable (no ACK)







# IEEE 802.11: DCF

- Why do we still get collisions with Collision Avoidance?
  - CA mitigates collisions, but it does not eliminate them
  - The collision probability depends on
    - the number of contending stations
    - the size of the contention window
  - A higher CW reduces the collision probability, but increases the delay introduced by back-off





# IEEE 802.11: DCF

- The DCF is:
  - A very strong MAC protocol in the context of a WLAN (an AP and 10-20 stations)
  - The most successful wireless technology in the world
  - The cheapest wireless technology on the market





# IEEE 802.11: DCF

- But the DCF is too:
  - Only a MAC layer solution among others (cellular networks use CDMA or OFDMA)
  - A mediocre protocol when mobility is considered
  - An unusable technology under high node density





# IEEE 802.11: DCF

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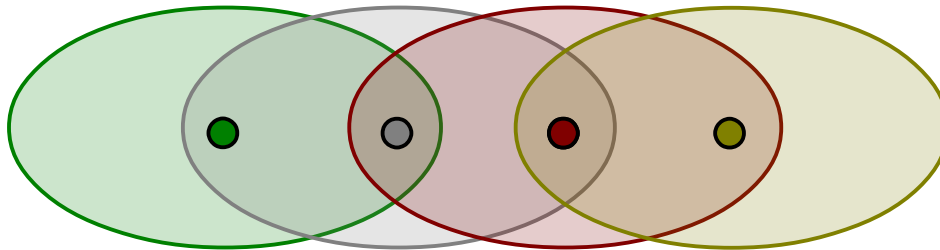




# IEEE 802.11: DCF

## An example – multi-hop networks

- **exposed terminal** – a tremendous reduction in throughput



- sends RTS to ●
- sends CTS to ●
- receives RTS from ● and refrain from transmitting to ●
- ...but transmission from ● to ● would not cause a collision!

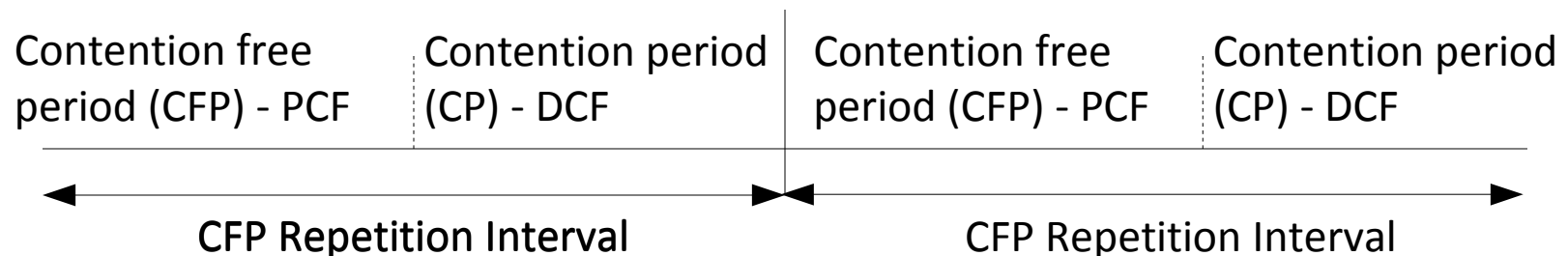






# IEEE 802.11: PCF

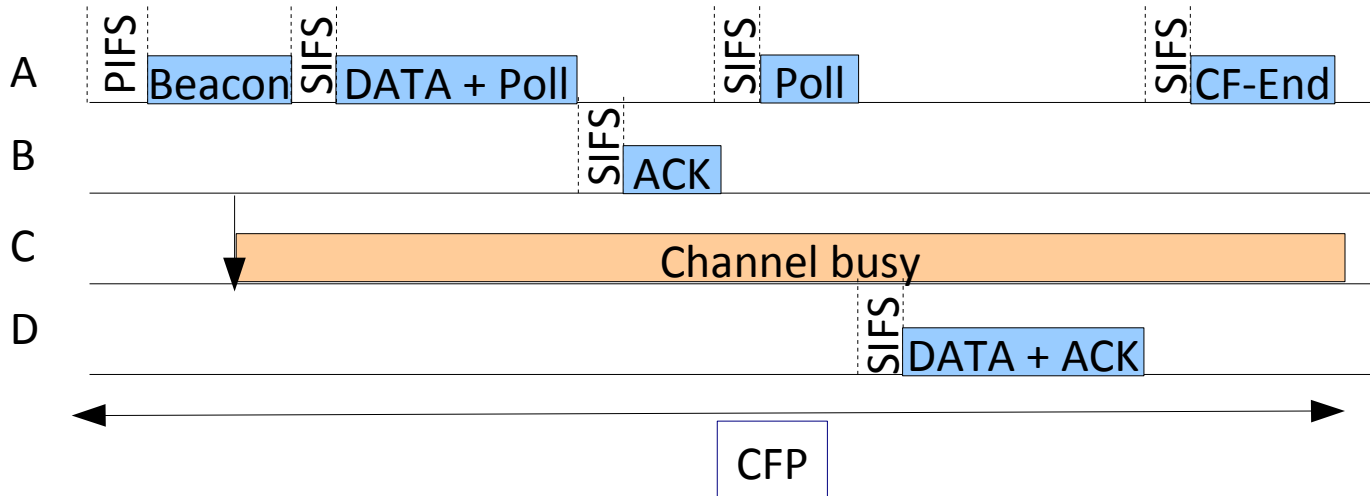
- There is more than DCF in IEEE 802.11
  - Point Coordination Function (PCF)
  - Contention-Free frame transfer protocol
  - Based on polling made by the access point
  - Coexists with DCF





# IEEE 802.11: PCF

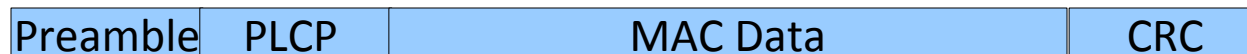
- How does PCF work?



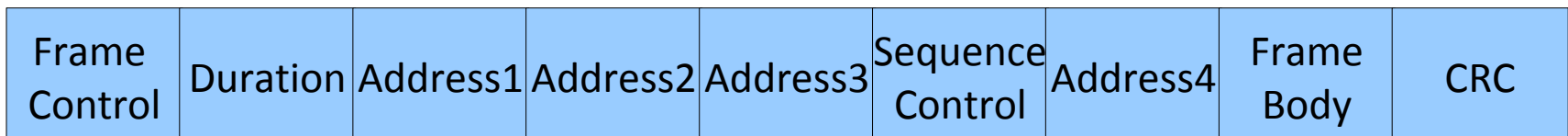


# IEEE 802.11: Frame Format

- The 802.11 frame



- The MAC Data





# IEEE 802.11: Frame Format

- Why do we need 4 addresses?
  - The Frame Control field contains (among others) two bits named To DS and From DS
  - The value of To DS and From DS gives the meaning of the 4 addresses

To DS	From DS	Address1	Address2	Address3	Address4
0	0	Destination	Source	BSSID	N/A
0	1	Destination	BSSID	Source	N/A
1	0	BSSID	Source	Destination	N/A
1	1	Receiver	Transmitter	Destination	Source





# CSMA/CA and IEEE 802.11

## Bibliography

- **IEEE 802.11 Working Group**, “*802.11: Wireless LAN Medium Access Control (MAC) and Physical Layer (PHY) Specifications*”, IEEE, 2007
- **G. Bianchi**, “*Performance Analysis of the IEEE 802.11 Distributed Coordination Function*”, IEEE Journal on Selected Areas in Communication, March 2000
- **S. Xu, T. Saadawi**, “*Does the IEEE 802.11 MAC protocol work well in multi-hop wireless ad hoc networks?*”, IEEE Communications Magazine, June 2001
- **A. Colvin**, “*CSMA with Collision Avoidance*”, Computer Communications, October 1983

